

## Components of a computer system

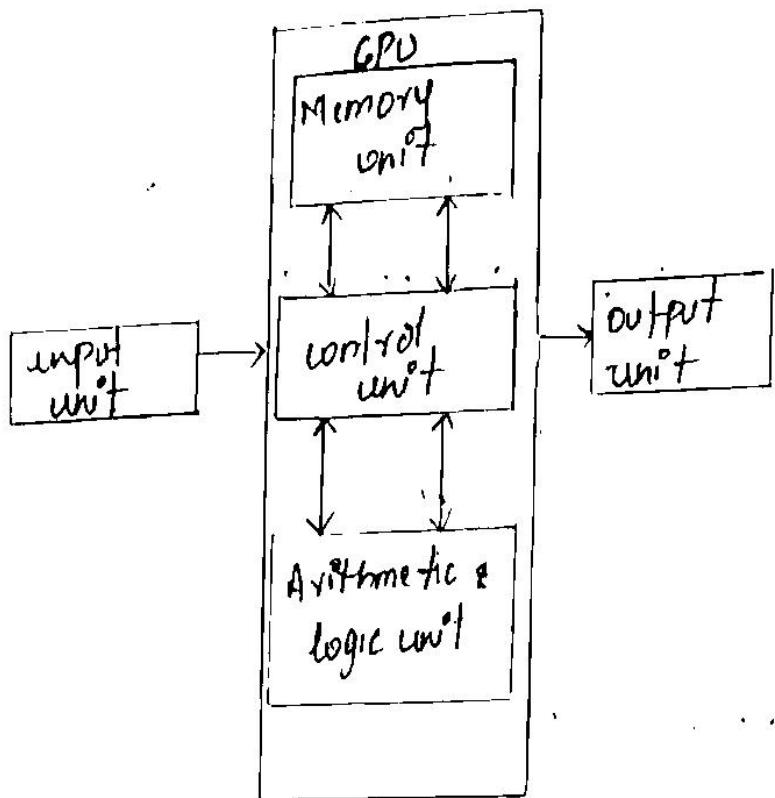
A computer is any machine that can be programmed to carry out a set of algorithms and arithmetic instructions.

A computer is a machine that can be instructed to carry out sequences of arithmetic or logical operations automatically via computer programming. Modern computers have the ability to follow generalized set of operations, called programs.

### Components of a Computer

All types of computers follow the same basic logical structure and perform the basic operations for converting raw data into information useful to their users.

- Motherboard
- central processing unit (CPU)
- Graphic processing unit (GPU) also known as video card.
- Random Access memory (RAM), also known as volatile memory.
- Storage: solid state drive (SSD) or Hard disk drive (HDD)



### Motherboard

Motherboard acts as a backbone of any technological device and it holds all the major components of a computer including the hard drive, processor, memory and peripheral ports like the USB.

Most motherboards are expandable i.e. we can replace components as long as they are compatible.

### Graphic card (or) video card

A graphics (or) video card can come in two varieties - Integrated (or) expansion.

An integrated video card is connected directly to the motherboard and is usually found as part of the processor.

An expansion video card is a separate card that is connected to another part of the motherboard called an expansion port.

The job of the video card is to create the graphics and images that can be shown on a monitor.

## Processor

The processor - also known as central Processing unit (CPU) is the brain of the computer.

CPU performs all types of data processing operations  
It stores data, intermediate results and instructions  
It controls the operation of all parts of the computer

CPU itself has the following three components

- ALU (Arithmetic logic unit)
- Memory unit
- control unit

### Arithmetic logic unit

This unit consists of two subsections namely

- Arithmetic section: is to perform arithmetic operations like addition, subtraction, multiplication and division.

Logic section: Function of logic section is to perform logic operations such as comparing, selecting, matching and merging of data.

### Memory (or) storage unit

This unit can store instructions, data and intermediate results. This unit supplies information to other units of the computer when needed. It is also known as internal storage unit or the main memory.

or the primary storage (a) Random Access Memory (RAM)

The size affects speed, power and capability.

Primary and secondary memory are two types of memories in the computer

Functions of the memory unit are

- It stores all the data and the instructions required for processing
- It stores ~~result~~ intermediate results for processing
- It stores the final results of processing before these results are released to an output device
- All inputs and outputs are transmitted through main memory

### Control Unit

This unit controls the operations of all parts of the computer but does not carry out any actual data processing operations.

Functions of the unit are

- It is responsible for controlling the transfer of data and instructions among other units of a computer
- It manages and coordinates all the units of a computer
- It obtains the instructions from the memory, interprets them, and directs the operation of a computer
- It communicates with input/output devices for transfer of data or results from storage
- It does not process (or) store data.

## Random Access Memory (RAM)

It is the internal memory of the CPU for storing data, program and rung result. It is a read/write memory which stores data until the system is working.

Access time in RAM is independent of the address, that is each storage location inside the memory is as easy to reach as other locations and takes the same amount of time.

Data in the RAM can be accessed randomly but it is very expensive.

Data is volatile i.e. data is erased (or) lost when we switch off the computer or on power loss.

RAM is of two types

- Static RAM (SRAM)
- Dynamic RAM (DRAM)

Hard drive : Hard drive is where programs and files are stored. HDD are comprised of a series of magnetized disks which store the data. These disks spin under a magnetic arm which can read and write data.

Modern hard drives are called solid-state drives (SSD) and use electrical circuits to store data. These are much faster than HDDs.

Static RAM - is one which requires the constant flow of the power to retain data inside it.  
capacitors

Dynamic RAM - needs to be refreshed to retain the data of hold-capacitors.

### ROM (Read only Memory)

- Rom - is non-volatile memory
- the capacity of ROM is comparatively smaller than RAM
- it is slower and cheaper than RAM
- the data in ROM can only be read by CPU
- ROM stores the instruction that computer requires during bootstrapping.

### Types of ROM

PROM - Programmable ROM - it can be modified only once by user

EPROM - electrically erasable and PROM - the content of the ROM can be erased using ultraviolet rays and the ROM can be programmed

EPRROM - electrically erasable PROM, it can be erased electrically and reprogrammed about ten thousand times

\* SSD are faster than HDD

## Introduction to Flowchart and Algorithms.

A set of steps involving arithmetic and computation and / or logical manipulation is required to obtain the desired result for a given problem. The desired result can be achieved through different courses of action and a variety of paths, and this is called "law of equifinality".

The most effective way of achieving a solution is called the optimum way. The optimum way of solving a problem to get the desired result can be achieved by analysing different strategies for the solution and then selecting the way that can yield the result in least time using the minimum amount of resources.

A type of analysis called task analysis is required to reach the solution from a problem definition that states what is to be achieved.

A set of steps that generates a finite sequence of elementary computational operations leading to the solution of a given problem is called an Algorithm.

Pictorial representation can be used as a substitute for an algorithm to understand more clearly and that is called as Flowchart.

Flowchart is a diagrammatic representation of an algorithm, boxes of different shapes are used to denote different types of operations and these boxes are connected by lines with arrows denoting the flow (or) direction. The connecting lines are known as flow lines.

Flowcharts may be classified into two categories:

i) Program flowchart

Prog flowchart are the mirror of computer program in terms of flowcharting symbols. They contain the steps of solving a problem unit for a specific result.

ii) System Flowchart

Contain the solution of many problem units together that are closely related to each other and interact with each other to achieve goal.

### Program flowchart

- Is an extremely useful tool in prog development
- Program flowchart is a pictorial representation of a logic of a program due to which it is easy to detect any error (or) omission from the program than from a program.
- Program flowchart can be followed easily and quickly.
- The type of documentation which may be of great help if the need for program modification arises in future.

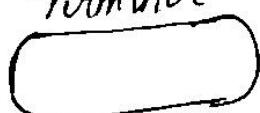
The following five rules should be followed while creating program flowcharts.

- 1) only the standard symbols should be used
- 2) the program logic should depict the flow from top to bottom and from left to right

- 3) each symbol used in a program flowchart should contain only one entry point and one exit point with the exception of decision symbol.  
This is known as single rule.
- 4) the operations shown within a symbol of a program flowchart should be expressed independently of any programming language.
- 5) All decision branches should be well-labeled.

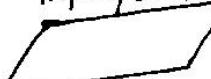
The following are the standard symbols used in program flowcharts.

Terminal



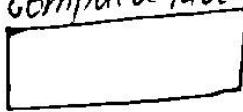
used to show the beginning and ending of a set computer-related processes

input/output



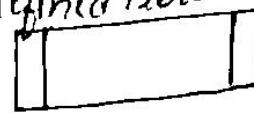
used to show any input/output operation

Computer Processing



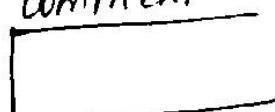
used to show any processing performed by a computer system

Predefined Processing



used to indicate any process not specially defined in the flowchart

Comment



used to write any explanation

Flow line

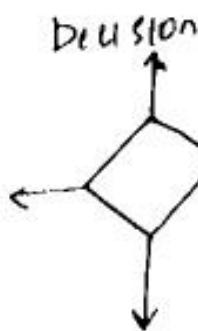


used to connect the symbols

Document I/O



used when input comes from the document and output goes to the document



Decision  
used to show any point in the process where a decision must be made to determine further action

on-page connector



used to connect parts of a flowchart continued on the same page

off-page connector



used to connect parts of a flowchart continued to separate pages

⇒ Flowcharts can be used to show the sequence of steps for doing any job.

⇒ A set of simple steps involving accepting inputs, performing arithmetic operation on the inputs, and showing them to the user demonstrate the sequence logic structure of a program

~~ess are faster~~

Flowchart

① The algorithm for the flowchart about cooking

step 1: Take the rice to be cooked

2: Procure the container

3: Procure the water

4: wash the rice in the water

5: put the rice into the container

6: Pour water into the container

7: if: water level = 1 inch above the rice

Then goto step 8

else goto step 6

endif

8: light the burner of the stove

9: if the rice is boiled

Then goto step 12

else goto step 10

10: endif

11: Heat the container

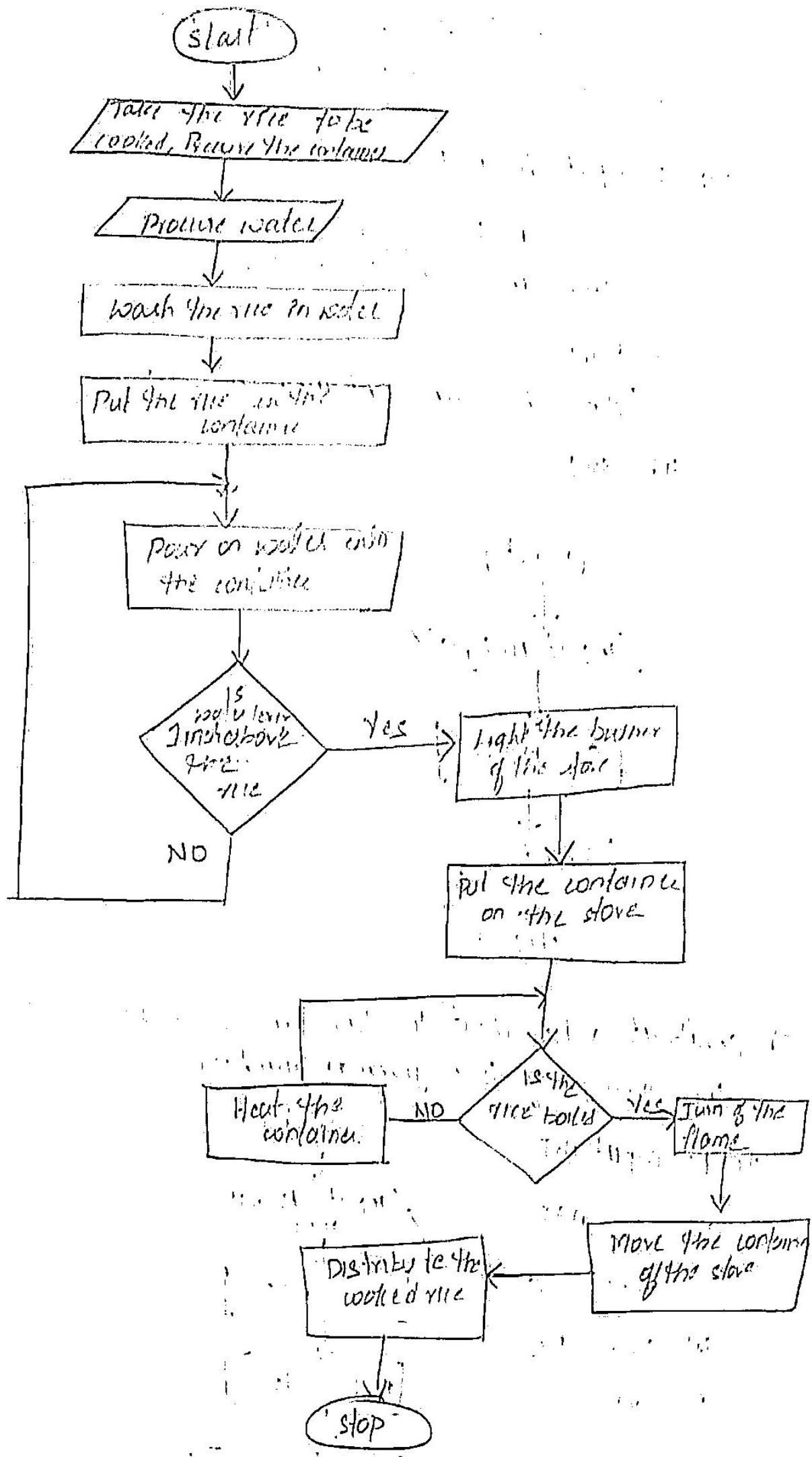
11: go to step 9

12: Turn off the flame

13: Move the container off the stove

14: Distribute the cooked rice

15: stop



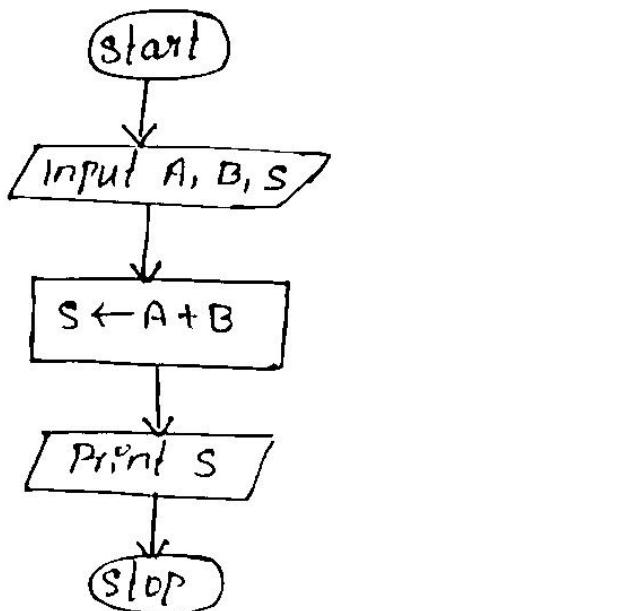
② Draw a flowchart to show how the sum of two numbers can be obtained.

step 1: Input to A, B

2:  $S \leftarrow A + B$   
(store the <sup>result</sup> of  $A + B$ )

3: Print S  
(show the sum obtained in step 2)

4: stop



③ construct a flowchart to show the procedure for obtain the average of two given numbers.

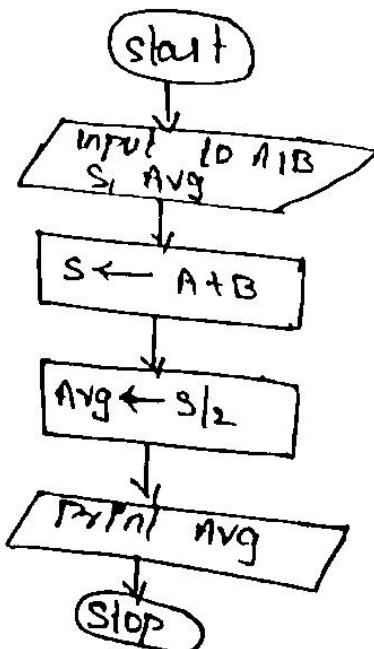
Step 1: Input A, B

2:  $S \leftarrow A + B$

3:  $Avg \leftarrow S/2$

4: Print Avg

5: stop



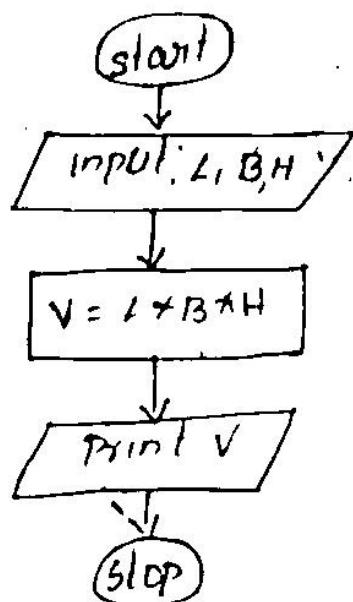
④ construct a flowchart to show how to obtain the volume of a rectangular box

step 1: input to L, B, H

2: compute  $V = L \times B \times H$

3: Print V

4: STOP



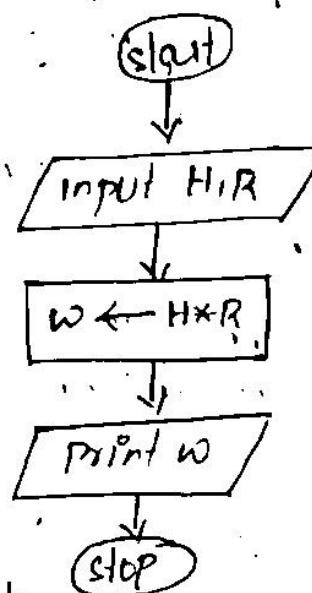
⑤ construct a flowchart to show how to obtain the daily wage of a worker on the basis of the hours worked during the day.

Step 1: input to H, R

2: compute  $w \leftarrow H \times R$

3: Print w

4: STOP



6) Find the product of two no's.

7) Find the remainder when one number is divided by the other

8) Find the area of a parallelogram

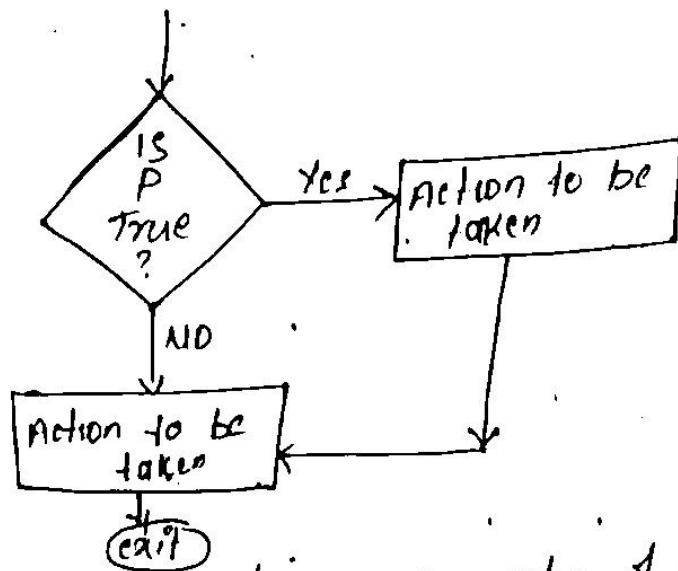
9) Find the area of a four-walls of a rectangular room

10) Find the area and perimeter of a circular plot

11) Find the area and perimeter of a square

## \* Decision-making

- the process of decision-making is implemented through a logic structure called selection
- here a predicate, also called a condition, is tested to see if it is true (or) false.



- A flowchart may contain any number of decision boxes depending on the processing requirements, and the boxes may appear in any sequence & depending on the program logic decided.

→ Develop a flowchart to show how the profit (or) loss for a sale can be obtained?

Algo Step 1: start

2: input to CP, SP

$SP > CP$

3: if ~~GP < SP~~       $50 < 40$

then goto step 4

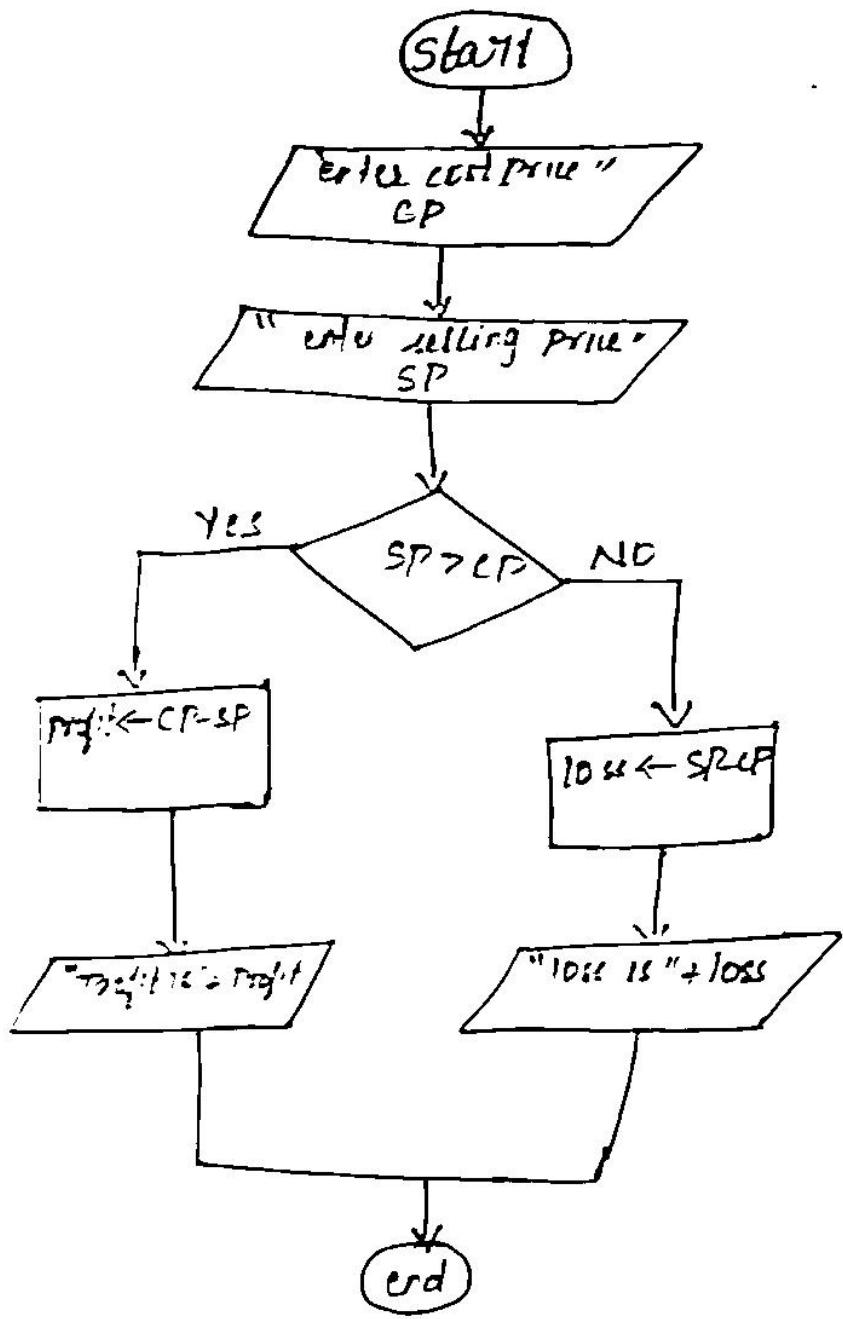
else goto step 5

4: end if

5:  $loss = SP - CP$

5:  $profit = CP - SP$

end



=, write an algo and draw flowchart to find the greatest  
of 3 no's

step 1: start

2: input to A, B, C

3 if  $A > B$  and  $A > C$

Print A

else go to step 4

4 if  $B > A$  and  $B > C$

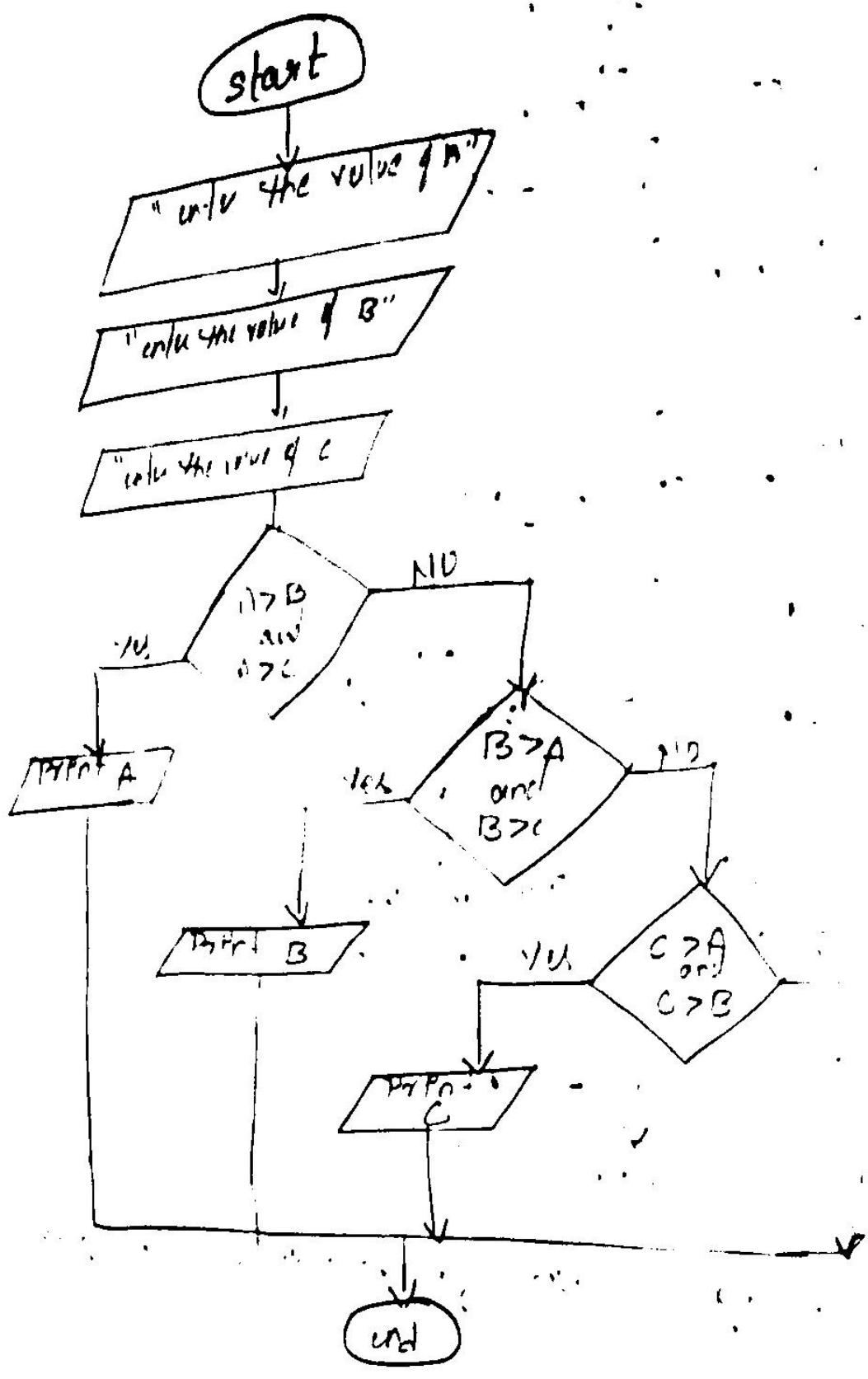
Print B

else go to step 5

5 if  $C > A$  and  $C > B$

Print C.

exit



⇒ write the algo and draw flowchart to find the given character is vowel or not.

## Fundamental Algorithms

### Exchanging the values of two variables

- Interchanging the values associated with two variables involves fundamental mechanism that occurs in sorting and data manipulation algorithms.

Suppose the variables  $a$  and  $b$  are assigned values

as      a      b  
      721      463

To exchange values of both the variables with the following assignments

$$\textcircled{1} \quad a := b \rightarrow \cancel{a = 721} \rightarrow 463 = b$$

$$\textcircled{2} \quad b := a \rightarrow 463 = a \times$$

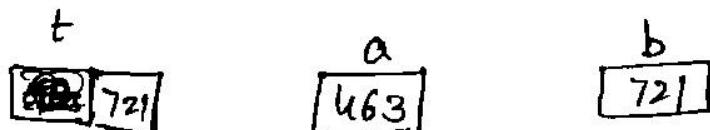
The first assignment would fine and exchanged the value in variable  $a$  but has left the value of  $b$  untouched.

To overcome this type of problem, need to introduce a temporary variable  $t$  and copy the original value of  $a$  into this variable before executing step  $\textcircled{1}$ .

$$t := a - \textcircled{3}$$

$$a := b - \textcircled{4}$$

$$b := t - \textcircled{5}$$



## Algorithm description

1. Save the original value of  $a$  in  $t$
2. Assign to  $a$  the original value of  $b$
3. Assign to  $b$  the original value of  $a$  stored in  $t$

## Pascal implementation

Procedure exchange (var a, b: integer)

Var var t: integer

begin {Save the original value of  $a$  then exchange  $a$  and  $b$ }

{ assert  $a=a \wedge b=b$  }

$t := a$

$a := b$

$b := t$

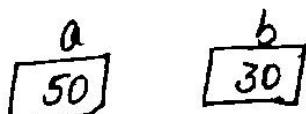
end.

use : The use of an intermediate temporary variable allows exchange of two variables to proceed correctly

## Application sorting algo.

- Q) Given two variables of integer type  $a$  and  $b$ , exchange their values without using a third temporary variable.

so



To exchange the values of both the variables with the following assignments

$$\text{Step 1: } a := b \rightarrow a := 30$$

$$2. b := a \rightarrow b := 30$$

value the step ① is successful in exchanging the value of variable a but failed in step ②

To exchange the values of variables without using third variable, can be done by using arithmetic operations.

i.e

$$\text{step 2} \quad \cancel{b = a} \quad \rightarrow \cancel{b = }$$

$$3. a := a + b \quad 50 + 30 = 80$$

$$4. b := a - b \quad 80 - 30 = 50$$

$$5. a := a - b. \quad 80 - 50 = 30$$

$$\begin{array}{c} a \\ \hline 30 \end{array} \quad \begin{array}{c} b \\ \hline 50 \end{array}$$

Algorithm description

- 1) add the values of both variables and save it in a
2. subtract the b value with a new value and assign to b
3. subtract the new value of b with new value of a and assign it to b.

## Counting

- counting mechanisms are very useful in computer algo
- count must be made of the number of items in a set which possess some particular property or which satisfy some particular condition or conditions

Suppose we are given the set of marks

55, 42, 71, 63, 29, 51, 89

with the condition  $\geq 50$

To make a count of this marks, let us start from left

first mark  $55 \geq 50 \rightarrow$  1 student marks is counted

Marks	Counting details for passed		
	Prev	Current	
55	0	1	$\geq 50 \checkmark$
42	1	1	$\geq 50 \times$
71	1	2	$\geq 50 \checkmark$
63	2	3	$\geq 50 \checkmark$
29	3	3	$\geq 50 \times$
51	3	3	$\geq 50 \times$
89	4	4	$\geq 50 \checkmark$
		5	$\geq 50 \checkmark$

The current count reflects the total no. of student passed the exam

To calculate the current count and previous-count

$$\text{so } \text{current\_count} = \text{Previous\_count} + 1 \quad -(1)$$

$$\text{Previous\_count} = \text{current\_count} \quad -(2)$$

these two steps must be repeatedly applied to obtain the count required.

The step(1) can also written as

$$\text{current\_count} = \text{current\_count} + 1$$

(New)

$$\text{ie } \text{old current\_count} = \text{previous\_count}$$

(or)

new current count as obtained by adding '1' to the old value of the current\_count and in such case the existence of previous\_count is unnecessary.

The essential steps in our pass-counting algorithm can be written as

while less than  $n$  marks have been examined do

(a) read next mark

(b) if current mark satisfies pass requirement then add one to count.

#### Algorithm description

1 prompt then read the number of marks to be processed

2 initialize count to 300

3 while there are still marks to be processed repeatedly do

(a) read next mark

(b) if it is a pass ( $e \geq 50$ ) then add count to count

4 write total number of passes.

## Pascal implementation.

Program passcount(input, output)

const passmark = 50;

var count & contains number of passes on termination

; { current number of marks processed by

m { current mark }

n { total number of marks to be processed first }

begin {count the number of passes (> 50) in a set of marks}

writeln('enter a number of marks n on a separate

line followed by the marks');

readln(n);

assert: n >= 0;

count := 0;

i := 0;

{invariant: count = number of marks in the first i  
read that are  $\geq$  passmark n  $\leq$  < n }

while i < n do

begin{read next mark, test it for pass and update  
count, if necessary}

i := i + 1;

read(m);

if coln(input) then readln;

if m  $\geq$  passmark then count := count + 1

end;

assert: count = number of passes in the set of n marks read

writeln('number of passes = ', count)

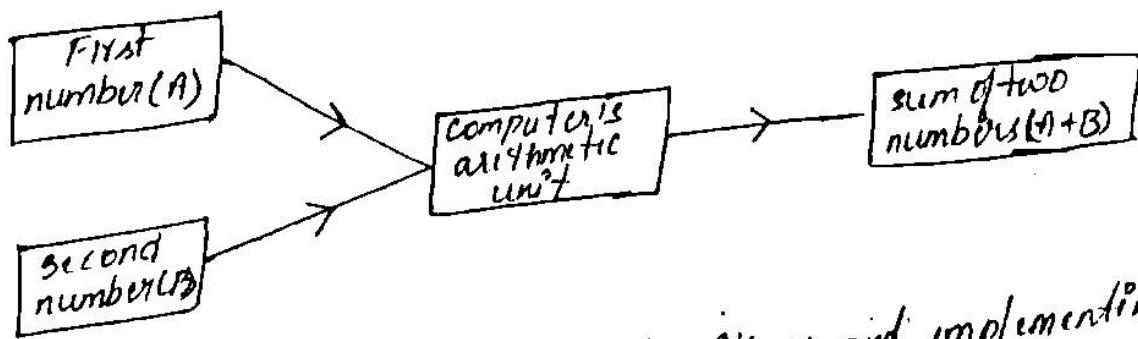
end.

- Initially, and each time through the loop, the variable count represents the number of passes.
- On termination, count represents the total number of passes in the set.
- Because 'i' is incremented by 1 with each iteration, eventually the condition  $i < n$  will be violated and the loop will terminate.

Applications:- All forms of counting

### Summation of a Set of Numbers

The computer has a built-in device which accepts two numbers to be added, performs the addition, and outputs sum of the two numbers.



- A fundamental goal in designing algorithms and implementing progs is to make the programs general enough to handle a wide variety of input conditions.

$$S = (a_1 + a_2 + a_3 + \dots + a_n)$$

or  $S = \sum_{i=1}^n a_i$  to add n numbers which is not practical

The core of the algorithm for summing  $n$  numbers involves a special step followed by a set of  $n$  iterative steps.

That is,

1. compute first sum ( $s=0$ ) as special case
2. Build each of the  $n$  remaining sums from the previous by an iterative process
3. Write out the sum of  $n$  numbers.

#### Algorithm description

1. Prompt and read in the number of numbers to be summed
2. Initialize sum for zero numbers
3. while does  $n$  numbers have been summed officially
  - (a) read in next number,
  - (b) compute current sum by adding the number read to the most recent sum .
4. write out sum of  $n$  numbers.

#### Pascal implementation

```
Program sum(input, output);  
Var i{summing loop index};  
n{number of numbers to be summed}: integer;  
a{current number to be summed};  
s{sum of n numbers on termination}: real;  
begin{computes sum of n real numbers for  
n>=0}
```

`writeln{'input n on a separate line followed by the,  
numbers to be summed'};`

`readln(n);`

`P := 0;`

`S := 0.0;`

{invariant :  $s = \text{sum. of first } P \text{ numbers readln}^i$ }  
`while i < n do`

`begin {calculate successive partial sums}`

`i := i + 1;`

`read(a);`

`if edn(input) then readln;`

`S := S + a;`

`end;`

`writeln{'sum of n = ', n, 'numbers = ', S};`

`end.`

### Notes on design

1. To sum  $n$  numbers  $(n-1)$  additions must be performed.
2. Initially, and each time through the loop, the sum  $s$  reflects the sum of first  $i$  numbers read.
3. On termination ( $i = n$ )  $s$  will represent the sum of  $n$  numbers.

### Applications

Average calculations, variance and least square calculations

- 2.3.1 Design an algo to compute the average of  $n$  numbers.
- 2.3.2 Redesign the algorithm so that it only needs to perform  $n-1$  additions to sum  $n$  numbers
- 2.3.3 Generate the first  $n$  terms of the sequence

1 2 4 8 16 32

### Factorial Computation

Given a number  $n$ , compute  $n$  factorial (written as  $n!$ ) where  $n \geq 0$

### Algorithm development

Start the development of algo by examining of  $n!$

$$n! = 1 \times 2 \times 3 \times \dots \times (n-1) \times n \text{ for } n \geq 1$$

and by the rule

$$0! = 1$$

$n!$  can always be obtained from  $(n-1)!$  by simply multiplying it by  $n$  (for  $n \geq 1$ ) that is

$$n! = n \times (n-1)! \text{ for } n \geq 1$$

let say  $P=0=1$  we can rewrite the steps as

$$P := 1 \quad -(1) = 0!$$

$$P := P \times 1 \quad = 1!$$

$$\left. \begin{array}{l} P := P \times 2 \\ P := P \times 3 \end{array} \right\} \rightarrow 2 \times 1 = 2! \quad = 3!$$

For the general  $(k+1)^{th}$  step we have ..

$$P_{k+1} = P_k \cdot k$$

- This general step can be placed in a loop to iteratively generate  $n!$ .
- This step allows us to take advantage of the fact that the computer's arithmetic unit can only multiply two numbers at a time.
- The central part of the algo for computing  $n!$  involves special initial step followed by  $n$  iterative steps.
  1. Treat 0! as a special case ( $P_0 = 1$ )
  2. Build each of the  $n$  remaining products  $P_i$  from its predecessor by an iterative process
  3. write out the value of  $n$  factorial

#### Algo description

1. Establish  $n$ , the factorial required where  $n \geq 0$
2. set product  $P$  for  $0!$ . Also set product count to 0.
3. while less than  $n$  products have been calculated  
    repeatedly do
  - (a) increment product count
  - (b) compute the  $i^{th}$  product if  $P$  by multiplying it by the most recent product.
4. return the result  $n!$

## Pascal implementation

```
function n!factorial(n: integer): integer;
var i {loop index representing ith factorial}: integer;
    factor {i!}: integer;
begin {computes and returns n! for n >= 0}
{assert: n >= 0}
    factors = 1;
    for i := 1 to n do
        factor := i * factor;
        n!factorial := factor;
end;
```

## Notes

This algorithm uses  $n$  multiplications to compute  $n!$

## Applications

Probability, statistical and mathematical computations

## Generation of the Fibonacci sequence.

Generate and print the first  $n$  terms of the fibonacci series where  $n \geq 1$

The first few terms are

0, 1, 1, 2, 3, 5, 8, 13.

Each term beyond the first two is derived from the sum of its two nearest preceding predecessors.

new term = preceding term + term before preceding term

### Algorithm description

1. Prompt and read  $n$ , the number of fibonacci numbers to be generated.
2. Assign first two fibonacci numbers  $a$  and  $b$ .
3. Initialize count of numbers generated.
4. While less than  $n$ :  
Fibonacci numbers have been generated do
  - (a) write out next two Fibonacci numbers;
  - (b) generate next fibonacci number keeping  $a$  relevant;
  - (c) generate next fibonacci number from most recent pair keeping  $b$  relevant for next computation
  - (d) update count of number of fibonacci numbers generated;
5. If  $n$  even then write out last two fibonacci numbers  
else write out second last fibonacci number

## Pascal Implementation

```
Program Fibonacci(input, output);
var a {fibonacci number variable};
     b {fibonacci number variable};
     i {number of fibonacci numbers generated};
     n {number of fibonacci numbers to be generated} : integer;
begin{block}[generate each fibonacci number from the sum of its two predecessors]
    a := 0;
    b := 1;
    i := 2;
    writeln('Enter n - the number of fibonacci numbers to be generated');
    readln(n);
    while i < n do
        begin
            writeln(a, b);
            a := a + b;
            b := a + b;
            i := i + 2;
        end;
    if i = n then writeln(a, b) else writeln(n-a);
end.
```

- Notes
- To generate  $n$  fibonaci numbers essentially  $n$  steps are required. The algo functions correctly for all values of  $n \geq 1$ .
  - Throughout the computation the variables  $a$  and  $b$  always contain the two most recently generated fibonaci numbers.
  - Applications.

The fibonaci sequence has numerous applications in biology, electrical network theory, counting and marking.

- Reversing the digits of an Integer
- Design an algorithm that accepts a positive integer and arranges the order of its digits.
- Digit reversal is a technique that is sometimes used in computing to remove bias from a set of numbers.
- It is important in fast information-retrieval algorithms.

### Algorithm description

1. Establish  $n$ , the positive integer to be reversed.
2. Set the initial condition for the reversed integer divisor.
3. While the integer being reversed is greater than zero do
  - (a) Use the remainder to extract the rightmost digit of the number being reversed
  - (b) Increase the previous reversed integer representation divisor by a factor of 10 and add it to the most recently extracted digit to give the current divisor value;

(c) use integer division by 10 to remove the rightmost digit from the number being reversed.

### Pascal implementation

```
function doreverse(n: integer): integer;
var reverse: integer;
begin{comment}reverse the order of the digits of a positive integer
  while n > 0 do
    begin
      reverse := reverse * 10 + n mod 10
      n := n div 10;
    end;
  end;
```

### Notes

- The number of steps to reverse the digits is directly proportional to the number of digits in the integer

### Applications

Hashing, and information retrieval, data base applications.

### Exp example

- design an algo that counts the number of digits in ~~an~~<sup>an</sup> integer
- design an algo to sum the digits in an integer